Algorithms engineering - Tommy Odland - Last edit: June 2, 2018

>Practial programming and Python

- Some modules from the *Python standard library* are indispensable in programming contests:
 - bisect for working with binary search.
 - defaultdict and deque from collections.
 - heapq for working with priority queues.
 - itertools for permutations, combinations, etc.
 - reduce from functools is sometimes useful.
 - See also math, string, operator and random.
 - Learn how to use the built-in data structures.
- Implementing a naive algorithm and testing against it on random inputs is often wise.
- If Python 3 is too slow, try Python 2. Use the %timeit and %lprun IPython functions to time code and identify slow parts. Use Java if all fails.

⊳Sliding, searching and sorting

- Comparison based sorting runs in $\mathcal{O}(n \log n)$ time.
- Maintaining a sorted list can speed up algorithms, since insertion is $\mathcal{O}(\log n)$ and retrieval is $\mathcal{O}(\log n)$. This advice has broad applicability.
- For problems "find minimal k such that P(k)", it might be possible to use binary search to find k if "Is P(k) possible?" can be answered quickly.

>Graph algorithms

- BFS and DFS are used to explore graphs. Both run in $\mathcal{O}(V + E)$ time. DFS uses a stack, BFS uses a queue. May be implemented recursively too.
- A *minimal spanning tree* may be found quickly using Krusal's algorithm or Prim's algorithm.
 - Prim's algorithm uses a priority queue.
 - Kruskal's algorithm uses a *disjoint-set data* structure to check if vertices are in the same set.
- Dijkstra's algorithm finds the single source shortest path to every other vertex in a graph.
 - A more general algorithm is the A^* search.

Dynamic programming

- Dynamic programming is used for problems with overlapping sub-problems, i.e. problems that obey the *principle of optimality*.
- The recipe is to (1) find a recursive formula, (2) identify the base cases (initial conditions) and (3) solve the problems such that no problem is solved twice. Either bottom-up to top-down using cache.

• Typically one creates a *DP table*, an *n*-dimensional table of function values. Dynamic programming is also possible over other data structures, such as trees and more general graphs.

>Exponential time algorithms

- Some algorithms run in exponential time, such as graph coloring.
- Approximate algorithms may often be used if finding the best answer is not necessary.
- If the exact answer is needed, use *backtracking* in an implicit tree to prune infeasible solutions as soon as possible to reduce run time in practice.

⊳Computational geometry

- The area of a *convex polygon* may be found in $\mathcal{O}(n)$ time using the shoelace formula.
- Given a set of n points in \mathbb{R}^2 , their convex polygon may be found in $\mathcal{O}(n \log n)$ time using *Graham's scan*, which orders the point by angle and examines them in a linear fashion using a stack.
- Angles, line crossings and so forth may be tricky due to floating point arithmetic and special cases, but typically relatively straight forward.

⊳Number theory

- The Sieve of Eratosthenes generates primes $\leq n$ in $\mathcal{O}(n \log \log n)$ time, since $1/2 + 1/3 + 1/5 + \cdots + 1/n = \log \log n$. It naively requires $\mathcal{O}(n)$ space.
- To verify if a number is prime, one option is to check numbers from 1 to \sqrt{n} , or use a precomputed prime list if available.

⊳Segment trees

- Given an array A, computing sums $A[i \dots j]$ has $\mathcal{O}(1)$ update and $\mathcal{O}(n)$ query. A cumulative sum has $\mathcal{O}(n)$ update and $\mathcal{O}(1)$ query.
- A segment tree balances the requirements above for $\mathcal{O}(n \log n)$ update and $\mathcal{O}(n \log n)$ query. It may also be used for other functions apart from sums.

⊳References

- Halim, Steven. *Competitive Programming*, 3rd Edition, 2013.
- Dasgupta, Sanjoy. Algorithms. McGraw Hill, 2008.